

## Announcements

Midterm 1 Wed in-class (50 minutes)

Reference sheet allowed (one A4 sheet, both sides)

Sections covered: 1.1-1.4, 2.1-3, 3.1-2, 5.1-2, 6.1 (no 6.2)

Practice problems partial solns posted

See policy email for more

Today: review

After class: problem session (review)

Tomorrow 10:30-11:30 am: office hour (via zoom @ snow day link)

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## Midterm 1 review:

### (Partial) list of topics:

Logic

Propositions

English vs. symbols

and, or, not, implies, for all, there exists

Truth tables

# Sets

Roster notation vs. set builder notation

Special sets ( $\mathbb{Z}$ ,  $\mathbb{R}$ ,  $\emptyset$ , etc.)

Venn diagrams

Subset, power set, Cartesian product

Cardinality

Set operations: union, intersection, set-minus, complement

Set identities (1-10)

Proof techniques: ett. chasing, membership tables

# Functions

Definition

Domain, codomain, range/image, preimage

Injective/surjective/bijective (b pf. techniques)

Composition

Inverses + invertibility

# Algorithms

Definition

Properties (describe and check)

Perform an algorithm

Write an algorithm

Searching/sorting/greedy change

## Big-O

Precise def'n of  $O$ ,  $\Omega$ ,  $\Theta$ ; proof techniques

Tricks & heuristics ( $1 < \log x < x < x^2 < \dots < e^x < \dots$ )

## Induction

Mathematical vs. strong

Base case, inductive step

Critique proofs

Various examples from class & H/W

## Counting

Sum/product/subtraction/division rules

Combining the rules (examples from class, HW)

## Examples:

1) Prove that  $A \cup (A \cap B) = A$  for all sets  $A, B$

Pf: We prove the result by showing that

a)  $A \cup (A \cap B) \subseteq A$  and b)  $A \subseteq A \cup (A \cap B)$ .

a) Let  $x \in A \cup (A \cap B)$ . Then either  $x \in A$  or  $x \in A \cap B$ .

In the former case, clearly  $x \in A$ , and in the latter case,  $x \in A$  and  $x \in B$ , so  $x \in A$ . Thus,  $A \cup (A \cap B) \subseteq A$ .

b) Let  $x \in A$ . Then  $x \in A$ , so  $x \in A \cup (A \cap B)$ . Thus,  $A \subseteq A \cup (A \cap B)$ . □

2) How many license plates can be made using either  
a) three digits followed by three uppercase letters or  
b) three uppercase letters followed by three digits

a) We have 10 choices for each of digits 1, 2, 3  
and 26 choices for each of digits 4, 5, 6, so by the  
product rule, there are

$$10 \cdot 10 \cdot 10 \cdot 26 \cdot 26 \cdot 26 = 17576000 \text{ such license plates}$$

b) We have 26 choices for each of digits 1, 2, 3  
and 10 choices for each of digits 4, 5, 6, so by the  
product rule, there are

$$26 \cdot 26 \cdot 26 \cdot 10 \cdot 10 \cdot 10 = 17576000 \text{ such license plates}$$

There is no overlap between the two cases, so in total  
there are

$$17576000 + 17576000 = 35152000 \text{ valid license plates}$$

3) Consider  $f: \mathbb{Z} \rightarrow \mathbb{Z}_{\geq 0}$  where  $f(x) = |x| + 1$ . Determine, with proof, whether  $f$  is inj., surj., bij., or none.

We claim that  $f$  is neither injective nor surjective.

Pf: First we consider injectivity.  $f$  is injective if and only if whenever  $f(x) = f(y)$ ,  $x = y$ . However,  $f(1) = 2 = f(-1)$ , so  $f$  is not injective.

Now,  $f$  is surjective if and only if  $f(\mathbb{Z}) = \mathbb{Z}_{\geq 0}$ ; in other words, if and only if for every  $y \in \mathbb{Z}_{\geq 0}$ , there exists  $x \in \mathbb{Z}$  s.t.  $f(x) = y$ . We claim that  $0 \notin f(\mathbb{Z})$ . To see this, notice that for all  $x \in \mathbb{Z}$ ,  $f(x) = |x| + 1 \geq 0 + 1 = 1 > 0$ . So  $0 \notin f(\mathbb{Z})$ , and  $f$  is not surjective.  $\square$

4) Name 3 of the 7 properties that algorithms should have, give a short description, and write some pseudocode which fails at only this property.

Properties: input, output, definiteness, correctness, finiteness, effectiveness, generality

**Finiteness:** For any input, the algorithm should produce the desired output after a finite number of steps

```
int-sqrt(pos integer x): (find  $\sqrt{x}$  if  $\sqrt{x} \in \mathbb{Z}$ ;  
   $i := 0$  otherwise, return -1)  
  while( $i^2 \neq x$ )  
     $i := i + 1$   
  return  $i$ 
```

**Correctness:** The algorithm should produce the correct output to the desired problem

```
int-sqrt(pos integer x):  
   $i := 0$   
  while( $x^2 \neq i$ )  
     $i := i + 1$   
  return  $i$ 
```

**Effectiveness:** It must be possible to perform each step exactly, and in a finite amount of time

```
int -sqrt (pos. integer x):  
   $A := \{(i, i^2) \mid i \in \mathbb{Z}_+\}$   
  If  $(i, x) \in A$  for some  $i \in \mathbb{Z}$   
    return  $i$   
  otherwise, return -1
```

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5) Find and prove a simple big- $\Theta$  estimate for  
$$f(x) = (7n^n + n2^n + 3^n)(n! + 3^n)$$

(Heuristic: Choose fastest-growing term of each factor)  
Let  $g(x) = n^n \cdot n!$ . We claim that  $f(x)$  is  $\Theta(g(x))$ .

Pf: We show that  $f$  is  $O(g)$  and  $g$  is  $O(f)$

Let  $C = 18$   $k = 10$

Then  $n! \geq 3^n$  for all  $x > k$  (this needs pf)

and  $n^n \geq n2^n$ ,  $n^n \geq 3^n$  for all  $x > k$  (this needs pf)

So for all  $x > k$ ,

$$\begin{aligned} |f(x)| &= (7n^n + n2^n + 3^n)(n! + 3^n) \\ &\leq (7n^n + n^n + n^n)(n! + n!) \\ &= 18n^n n! = C|g(x)|, \end{aligned}$$

So  $f$  is  $O(g)$ .

(Other direction:  $C = k = 1$  works)